Precept Outline

- Review of Lectures 5 and 6:
 - Comparators and Comparables
 - Elementary sorts
 - Mergesort

A. Review: O/Ω Notation + Elementary Sorts + Mergesort + Comparable/Comparator

Your preceptor will briefly review key points of this week's lectures. They may refer to the warm-up exercise and the code snippet shown below.

Warm-up: Let $f(n) = 3n + 4n \log_2 n + 8\sqrt{n} \log_2 n$. Select all that apply.

 $\begin{array}{l} (\ \) \ f(n) = O(n) \\ (\ \) \ f(n) = \Omega(n) \\ (\ \) \ f(n) = \Omega(n) \\ (\ \) \ f(n) = O(\sqrt{n}\log n) \\ (\ \) \ f(n) = \Omega(n\log n) \\ (\ \) \ f(n) = \Omega(n\log n) \\ (\ \) \ f(n) = O(n^2) \\ (\ \) \ f(n) = \Omega(n^2) \\ (\ \) \ f(n) = \Omega(n^2) \\ (\ \) \ f(n) = \Omega(\log n) \\ (\ \) \ f(n) = O(\log n) \\ (\ \) \ f(n) = O(\log n) \\ (\ \) \ f(n) = O(\log n) \\ (\ \) \ f(n) = O(\log n) \\ (\ \) \ f(n) = O(2^n) \\ \end{array}$

() $f(n) = \Omega(2^n)$

```
public class YourClass implements Comparable<YourClass> {
     public int compareTo(YourClass that) {
2
          // returns int > 0 if this > that
3
          // returns int < 0 if this < that</pre>
4
          // returns 0 otherwise
     }
     private static class YourComparator implements Comparator<YourClass> {
              public int compare(YourClass obj1, YourClass obj2) {
              // returns int > 0 if obj1 > obj2
              // returns int < 0 if obj1 < obj2</pre>
              // returns 0 otherwise
2
          }
3
4
     }
     public static Comparator<YourClass> yourComparison() {
5
          return new YourComparator();
6
7
     }
8
      . . .
9 }
```

- **Relevant Book Sections**
- Book chapters: 2.1, 2.2 and 2.5

B. Comparable & Comparator

The code snippet below shows the instance variables of a class Movie, and partially filled instance methods that should support comparing elements of this class in three ways:

- by alphabetical order of title (the default order);
- by release year; and
- by rating (0-5 stars).

Fill in the blanks numbered 1 to 6.

```
public class Movie implements _____(1)_____ {
1
     private String title;
2
     private int year;
3
     private int rating;
5
     public int compareTo(Movie m) {
6
        return _____(2)_____;
     }
8
9
0
     public static Comparator<Movie> byYear() {
1
        return new YearComparator();
2
     }
3
     private static class YearComparator implements _____(3)_____ {
4
        public int compare(Movie m1, Movie m2) {
5
6
            return _____(4) _____;
7
        }
8
     }
19
     public static Comparator<Movie> byRating() {
20
       return new RatingComparator();
     }
23
24
25
     private static class RatingComparator implements _____(5)_____ {
        public int compare(Movie m1, Movie m2) {
           return _____(6)____;
        }
     }
29
     . . .
$0 }
```

C. Sorting Algorithms

Part 1: Spring'24 Midterm Problem

Given two integer arrays, a[] and b[], the *symmetric difference* between a[] and b[] is the set of elements that appear in exactly one of the arrays. Design an algorithm that receives two *sorted arrays*, each consisting of n *distinct elements*, and outputs the size of their symmetric difference.

For full credit, it must use $\Theta(1)$ extra memory and its running time must be $\Theta(n)$ in the worst case (the arrays a[] and b[] should not be modified). A solution with $O(n \log n)$ runtime and O(n) extra memory that does not satisfy the full credit performance requirements receives partial (at least half) credit.

Part 2: Sorting Lower Bounds

Imagine you are given unlimited access to call a method (say, via "the cloud") which costs your program *constant time* in order to help sort an array.

(a) Suppose the method is sum(int[] a, int i, int j), which, given two indices $0 \le i \le j < n$, returns the sum $\sum_{k=i}^{j} a[k]$. Can you use it to implement a (comparison-based) sorting algorithm with O(n) running time? If so, how? If not, why not?

(b) Suppose the method is min(int[] a, int i), which returns $\min_{i \le k < n} \{a[k]\}$. Can you use it to implement a (comparison-based) sorting algorithm with O(n) running time? If so, how? If not, why not?

Part 3: Equality of Histograms

The *histogram* of an array s[] of samples is the set of pairs (i, f_i) , where f_i is the number of indices j such that the j^{th} sample s[j] has value i. (That is, $f_i = |\{j : s[j] = i\}|$.)

Let a[] and b[] be integer arrays representing sample sequences. Design an algorithm with $O(n \log n)$ worst-case running time that identifies whether the histograms of a and b are equal (i.e., if, for all *i*, the frequency of *i* is the same in a and b).

D. Assignment Overview: Autocomplete

Your preceptor will introduce and give an overview of your third assignment. Don't hesitate to ask questions! Summary of the assignment:

- Implement a Term class, which stores a word (as a string) and a numeric weight, and also implements comparators for comparing terms in natural order, in decreasing order of weight, and lexicographically based on the first *r* characters.
- Create a data type Autocomplete that initializes with given arrays of terms and weights, and supports methods to return the weight of a term, the top matching term, and the top *k* matching terms in descending order of weight.
- Implement a BinarySearchDeluxe class, which should use binary search to find the first and last index of a given key in a sorted array (these are important primitives to the Autocomplete class).

E. Optional Bonus Problems

Part 1: Three-way Mergesort

(Two-way) Mergesort is quite a simple algorithm to describe: to sort n elements, divide the array in half, (recursively) sort each then merge the two halves together. In this exercise, we will study a variant of it: in three-way Mergesort, we divide an array of length n into 3 subarrays of length $\frac{n}{3}$, sort each of them and then perform a 3-way merge.

Given 3 **sorted** subarrays of size $\frac{n}{3}$, how many comparisons are needed (in the worst case) to merge them to a sorted array of size *n*? Provide your answer in tilde notation.

What is the order of growth of the number of compares in 3-way Mergesort as a function of the array size n? (Here we're counting the total number, including all recursive calls.)

Given a choice, would you choose 3-way or 2-way mergesort? Justify your answer.

Part 2: (Challenge) Counting Inversions

In an array h of n numbers, an *inversion* is a pair of elements that isn't sorted; that is, two indices i and j such that i < j and h[i] > h[j].

Describe an algorithm to compute the total number of inversions of an array of length n in time $\Theta(n \log n)$. *Hint: think about how you can modify mergesort to achieve this.*